# Some structural similarities between uncountable sets, powersets and the universe

Athanassios Tzouvaras

Department of Mathematics Aristotle University of Thessaloniki 541 24 Thessaloniki, Greece e-mail: tzouvara@math.auth.gr

#### Abstract

We establish some similarities/analogies between uncountable cardinals or powersets and the class V of all sets. They concern mainly the Boolean algebras  $\mathcal{P}(\kappa)$ , for a regular cardinal  $\kappa$ , and  $\mathcal{C}(V)$  (the class of subclasses of the universe V), endowed with some ideals, especially the ideal  $[\kappa]^{<\kappa}$  for  $\mathcal{P}(\kappa)$ , and the ideal of sets V for  $\mathcal{C}(V)$ .

Mathematics Subject Classification (2020): 03E05, 03E20.

*Keywords*: Uncountable cardinal, access ideal for a set/class, monadic second-order language of ordering.

### 1 Introduction

This article belongs roughly to the same line of thought as [6] and [8], which on the one hand try to shed some light to the structure of powersets of infinite sets, and on the other, at philosophical level, express a skeptical attitude towards the powerset axiom (Pow), that is the claim that the collection of subsets of an infinite set is a set (rather than a proper class). To the same vein belongs also [9], where it is proposed that the powerset and replacement axioms (especially the former one) would be better to be replaced by local (and thus relativized) versions of them that lead to a local version of the entire ZFC. The approach of the present article differs, however, in that we focus on the relationship and similarities between *uncountable* sets in general, rather than powersets alone, and the universe. Since the connection between the existence of powersets and the existence of uncountable sets seems somewhat confused, let me elaborate briefly on this point before coming to the main body of the paper.

It is well-known that powersets and uncountable sets are closely related through the famous Cantor's theorem that every infinite powerset is uncountable. However this connection seems to be only one-way: as is wellknown, the existence of uncountable sets is (relatively) consistent with the failure of *Pow*. For example if  $M \models \text{ZFC} + 2^{\aleph_0} = \aleph_2$ , then

$$H^{M}_{\omega_{2}} \models \text{ZFC}^{-} + ``\omega_{1} \text{ exists''} + ``\mathcal{P}(\omega) \text{ does not exist''},$$

where  $H_{\omega_2} = \{x : |TC(x)| < \aleph_2\}$  and  $ZFC^- = ZFC - \{Pow\}$ . So, if *Unc* denotes the existence of uncountable sets,  $ZFC^- + Unc \not\vdash Pow$ , i.e., *Pow* is *formally* independent of *Unc* over  $ZFC^-$ .

The problem is that a resident of  $H^M_{\omega_2}$ , in contrast to a resident of M, has no clue as to how the uncountable  $\omega_1$  has emerged. The resident of M has a natural way to construct  $\omega_1$ , as well as the other uncountable cardinals: they consider the collection  $WO(\omega)$  of all well-orderings of  $\omega$ , which is a subcollection of  $\mathcal{P}(\omega \times \omega)$ . Since the latter is a set by Pow, so is  $WO(\omega)$ by Separation. Then taking the set  $WO(\omega)/\cong$  of isomorphism types of  $WO(\omega)$ , they arrive at  $\omega_1$ . But for the resident of  $H^M_{\omega_2}$ ,  $\mathcal{P}(\omega \times \omega)$  is only a proper class. It is true that they can define  $WO(\omega)$  in  $H^M_{\omega_2}$ , as a subclass of  $\mathcal{P}(\omega \times \omega)$ . But they cannot even define  $WO(\omega)/\cong$ , since each isomorphism type of a well-ordering of  $\omega$  has continuum many members, so it is a proper class too. Thus  $\omega_1$  is a completely accidental object of  $H^M_{\omega_2}$ , which seems to be there by mere chance.

If, alternatively, we start with a model  $N \models \text{ZFC} + 2^{\aleph_0} = \aleph_1$ , then

$$H_{\omega_2}^N \models \operatorname{ZFC}^- + \mathscr{P}(\omega) \text{ exists}^* + \mathscr{P}(\mathcal{P}(\omega)) \text{ does not exist}^*,$$

and the problem now for a resident of  $H_{\omega_2}^N$  is to explain how the uncountable set  $\mathcal{P}(\omega)$  has occurred, since the operation  $\mathcal{P}(x)$  is not defined for all sets.

The role of Pow in ZF is quite analogous to the role of exponentiation axiom (Exp) in (subsystems) of Peano arithmetic (PA). Exp is the sentence (in the language of PA)  $\forall x, y \exists z (x^y = z)$ , or more simply  $\forall x \exists y (2^x = y)$ . If  $I\Sigma_n$  denotes the fragment of PA that consists of the elementary axioms of PA and the induction scheme restricted to  $\Sigma_n$  formulas alone, then it is wellknown that Exp is provable in  $I\Sigma_1$  but unprovable in  $I\Sigma_0$ . More precisely,  $I\Sigma_0 < I\Sigma_0 + Exp < I\Sigma_1$  (where for theories  $T_1, T_2, T_1 < T_2$  means that  $T_2$  is strictly stronger than  $T_1$ ). Roughly, Exp is with respect to  $I\Sigma_0$ , what Pow is with respect to ZF<sup>-</sup>, and uncountable sets are analogous to "very big" numbers of arithmetic. The only difference is that uncountable sets have a precise definition (sets of size strictly larger than  $\aleph_0$ ), while very big numbers do not. Intuitively *a* is "very big" if  $a \geq 2^b$  for some "big" *b*. As is wellknown the models of I $\Sigma_0$  are closed under + and  $\cdot$  but not, in general, under exponentiation. In particular, the operation  $x^y$  is only partially defined in them. Nevertheless, there are (nonstandard) models of I $\Sigma_0 + \neg Exp$  (see e.g. [5], Theorems 4.3, 4.4) which, despite the failure of Exp, may contain  $2^a$  for some arbitrary nonstandard number *a*. In the absence of the full exponentiation operation, the existence of  $2^a$  can be seen by a resident of the model only as *accidental*, exactly as was seen the cardinal  $\omega_1$  in the model  $H_{\omega_2}^M$  of ZFC<sup>-</sup> mentioned above.

The conclusion of the preceding discussion is that the connection between powersets and the uncountable is not just one-way. It is in fact a two-way dependence, if behind every encounter with an uncountable set we seek the reason and source of its occurrence.

Having argued that Pow is, directly or indirectly, responsible for the existence of uncountable sets, a reasonable way to reinforce the skeptical attitude towards this axiom is to present concrete mathematical facts that establish specific similarities/analogies between uncountable sets and the universe of all sets V. This is what we are doing in this article. Namely, we compare the Boolean algebra  $\mathcal{P}(\kappa)$  of an infinite cardinal  $\kappa$ , with the Boolean algebra  $\mathcal{C}(V)$  of subclasses of V, and point out certain similarities. The similarities are of three kinds: (a) set-theoretic, (b) algebraic, and (c) order-theoretic. In the comparison of set-theoretic and algebraic facts some ideals play a key role, namely the ideals  $[\kappa]^{<\lambda}$ , for  $\lambda \leq \kappa$ , and in particular  $[\kappa]^{<\kappa}$ , for the case of cardinals, and the ideal of sets V for  $\mathcal{C}(V)$ . We deal with each of the three kinds of similarities in separate sections.

**Summary of contents.** In section 2 we deal with a specific set-theoretic similarity. The main results are the following:

(a) The ideal V of  $\mathcal{C}(V)$  is *minimal* in a natural and precisely defined way. Moreover, it is the *unique* minimal ideal of this kind (Theorem 2.5).

(b) For every regular cardinal  $\kappa$ , the ideal  $[\kappa]^{<\kappa}$  of  $\mathcal{P}(\kappa)$  is *minimal* in the above sense, and also the *unique* minimal ideal of this kind (Theorem 2.6).

(c) (In sharp contrast to (b)), for every singular cardinal  $\kappa$ , there is *no* minimal ideal of  $\mathcal{P}(\kappa)$  in the above sense (Theorem 2.11).

In section 3 we deal with the algebraic similarities. Specifically we compare the Boolean algebra  $\mathcal{P}(\kappa)$  equipped with an ideal  $[\kappa]^{<\lambda}$ , for any

 $\aleph_0 \leq \lambda \leq \kappa$ , with the Boolean algebra  $\mathcal{C}(V)$  equipped with the ideal V. Using Ehrenfeucht-Fraïssé games, it is shown that these structures are elementarily equivalent, that is,  $\langle \mathcal{P}(\kappa), [\kappa]^{<\lambda} \rangle \equiv_{L_b(S)} \langle \mathcal{C}(V), V \rangle$ , where  $L_b(S)$  is the language of Boolean algebras augmented with a unary predicate  $S(\cdot)$ .

In section 4 we deal with the order-theoretic similarities. For convenience we identify V with the class of ordinals On by the help of von Neumann's global choice vN, and compare the first-order structures  $\langle \kappa, < \rangle$  and  $\langle On, < \rangle$ , where < is the natural well-orderings of On and  $\kappa$  (i.e.,  $\in$ ). Extending results of [2] and [7] we show that  $\langle \kappa, < \rangle \equiv_{L_{ord}} \langle On, < \rangle$ , where  $L_{ord}$  is the first-order language of ordering. This equivalence is further strengthened to  $\langle \kappa, <, \operatorname{Cof}(\kappa) \rangle \equiv_{pos}^{\forall_1^1} \langle On, <, \operatorname{Cof}(On) \rangle$ , where  $\operatorname{Cof}(\kappa)$  (resp.  $\operatorname{Cof}(On)$ ) denotes the class of cofinal subsets of  $\langle \kappa, < \rangle$  (resp.  $\operatorname{Cof}(On)$ ) denotes the class of cofinal subsets of  $\langle \kappa, < \rangle$  (resp.  $\operatorname{Cof}(On)$ ) denotes the class of the monadic (second-order) language  $L_{mon}$ , which is  $L_{ord}$ augmented with set variables and  $\in$ .

Some notational conventions. In order to talk about subclasses of V we need to work in a class theory (rather than ZFC), like Gödel-Bernays theory with choice GBC. So throughout our discussion will be held having as ground theory either GBC or ZFC. Concerning the collection  $\mathcal{C}(V)$  of all subclasses of V when working in GBC, this is actually a "virtual", i.e. syntactic, entity, since in order to accommodate it literally as an extensional object, we would need a third-order set theory. However this is not a serious problem. We can always treat  $\mathcal{C}(V)$  in model-theoretic terms, that is, think that we work in an arbitrary model  $\mathfrak{M} = \langle M, \mathcal{M} \rangle$  of GBC, where M contains the sets and  $\mathcal{M}$  contains the classes of  $\mathfrak{M}$ . Then  $V^{\mathfrak{M}} = M$  while  $\mathcal{C}(V)^{\mathfrak{M}} = \mathcal{M}$ . Thus  $\mathcal{C}(V)$  is simply the range of class variables and, like  $\mathcal{P}(A)$ , an atomic Boolean algebra. On the other hand, when we talk about *ideals* of  $\mathcal{C}(V)$ , which are collections  $\mathcal{I} \subseteq \mathcal{C}(V)$ , these will always be *definable* by a formula  $\phi(X)$  with one free class variable.

We have already used above the notation  $[\kappa]^{<\lambda}$ . In general, for any set A and cardinal  $\kappa \leq |A|$  we let

$$[A]^{<\kappa} = \{x \subseteq A : |x| < \kappa\}^1, \quad [A]^{\leq \lambda} = \{x \subseteq A : |x| \leq \lambda\}.$$

Also, in view of AC, we can identify every powerset  $\mathcal{P}(A)$  with a cardinal  $2^{\kappa}$ , and write  $\mathcal{P}(2^{\kappa})$  instead of  $\mathcal{P}(\mathcal{P}(A))$ ,  $[2^{\kappa}]^{<\lambda}$  instead of  $[\mathcal{P}(A)]^{<\lambda}$ , and so on.

<sup>&</sup>lt;sup>1</sup>This set is usually denoted also  $\mathcal{P}_{\kappa}(A)$ .

# 2 A set-theoretic similarity between V and regular cardinals: minimality and uniqueness of access ideals

Throughout upper case letters X, Y are used as class variables for the language of GBC, while lower case letters x, y denote sets (elements of classes). For two classes X, Y, let |X| = |Y| denote the fact that X and Y are equipollent, i.e. there is a bijection between them. The axiom vN says that all proper classes are equipollent, or equivalently |V| = |On|. Following the conventions agreed in the Introduction we shall write  $2^{\aleph_0}$  instead of  $\mathcal{P}(\omega)$  and in general  $2^{\kappa}$  instead of  $\mathcal{P}(\kappa)$ .

Definable classes are produced as usual by formulas  $\phi(x)$  with a set variable, i.e.,  $X = \{x : \phi(x)\}$ . But also classes of classes are permitted, provided they are defined by formulas  $\phi(X)$  with a class free variable. We denote such classes by calligraphic letters  $\mathcal{I}$ ,  $\mathcal{J}$  and treat them as syntactic objects, namely  $\mathcal{I}$ ,  $\mathcal{J}$  can be thought of as formulas, and writing  $\mathcal{I} \subseteq \mathcal{J}$  we just mean that  $(\forall X)(\mathcal{I}(X) \to \mathcal{J}(X))$ . For any such  $\mathcal{I}$  we have  $\mathcal{I} \subseteq \mathcal{C}(V)$ , where  $\mathcal{C}(V)$  is the class of all subclasses of V. Given a class X, an *ideal* of  $\mathcal{C}(X)$  (or an ideal on X) is a class  $\mathcal{I} \subseteq \mathcal{C}(X)$  such that:

- (i) If  $Y_1, Y_2 \in \mathcal{I}$ , then  $Y_1 \cup Y_2 \in \mathcal{I}$ ,
- (ii) If  $Y \in \mathcal{I}$  and  $Z \subseteq Y$ , then  $Z \in \mathcal{I}$ ,
- (iii)  $\mathcal{I}$  is proper, i.e.,  $X \notin \mathcal{I}$ .

Trivially V is an ideal of  $\mathcal{C}(V)$ .

#### 2.1 Accessibility of classes and sets by ideals

A (total) preordering on a class X is a binary relation  $\preccurlyeq$  which is reflexive, transitive and total, where the latter means that for all  $x, y \in X$   $x \preccurlyeq y$ or  $y \preccurlyeq x$ .  $x \prec y$  means  $x \preccurlyeq y$  and  $x \neq y$ . If we add the antisymmetric property, the preordering becomes a linear ordering. To avoid trivialities we shall assume throughout that no preordered class  $\langle X, \preccurlyeq \rangle$  has a last element. Given a preordering  $\langle X, \preccurlyeq \rangle$  and  $x \in X, \preccurlyeq_x$  denotes the initial segment of X determined by x, i.e.,  $\preccurlyeq_x = \{y \in X : y \preccurlyeq x\}$ . A  $Y \subseteq X$  is said to be bounded w.r.t.  $\preccurlyeq$  if there is  $x \in X$  such that  $Y \subseteq \preccurlyeq_x$ . An unbounded  $Y \subseteq X$ is called also cofinal. For every preordering  $\preccurlyeq of X$ , the bounded subsets of X form a natural ideal of  $\mathcal{C}(X)$  denoted  $Seg(\preccurlyeq)$ , that is

$$Seg(\preccurlyeq) = \{Y \subset X : \exists x \in X (Y \subseteq \preccurlyeq_x)\}.$$

The fact that  $\langle X, \preccurlyeq \rangle$  has no last element guarantees that the ideal  $Seg(\preccurlyeq)$  is proper.

**Definition 2.1** (GBC or ZFC) Let  $\mathcal{I}$  be an ideal on a class or set X.

(1)  $\mathcal{I}$  is said to be an *access ideal for* X, or X *is*  $\mathcal{I}$ *-accessible*, if there is a preordering  $\preccurlyeq$  of X such that  $Seg(\preccurlyeq) \subseteq \mathcal{I}$ . In that case  $\preccurlyeq$  is said to be an  $\mathcal{I}$ *-preordering*.

(2) The ideal  $\mathcal{I}$  is said to be a minimal access ideal, or just minimal for X, if it is an access ideal for X and for any ideal  $\mathcal{J} \subsetneq \mathcal{I}, X$  is not  $\mathcal{J}$ -accessible.

Trivially every preordering  $\preccurlyeq$  of X is a  $Seg(\preccurlyeq)$ -preordering. Notice that, although the preceding definitions are expressible in GBC or ZFC (when Xand  $\mathcal{I}$  are definable objects), it is not to be meant that it is also *decidable* in these theories whether, for given X and  $\mathcal{I}$ , (a) X is  $\mathcal{I}$ -accessible or not, and (b)  $\mathcal{I}$  is minimal for X or not. These are relative facts, in general, depending on the model of GBC/ZFC we work in. For example, as we shall see below, the cardinal  $2^{\aleph_0}$  is accessible by the ideal  $[2^{\aleph_0}]^{\leq\aleph_0}$  of its countable subsets if and only if  $2^{\aleph_0} = \aleph_1$ . Moreover, in the latter case  $[2^{\aleph_0}]^{\leq\aleph_0}$  is the unique minimal access ideal for  $2^{\aleph_0}$ . More generally, if  $2^{\aleph_0} = \aleph_{\alpha}$  holds in a model  $\mathfrak{M}$  and  $\aleph_{\alpha}$  is regular, then  $[2^{\aleph_0}]^{<\aleph_{\alpha}}$  is the unique minimal access ideal for  $2^{\aleph_0}$  in  $\mathfrak{M}$ . In contrast, if  $\aleph_{\alpha}$  is singular, then  $2^{\aleph_0}$  is  $[2^{\aleph_0}]^{<\aleph_{\alpha}}$ -accessible but there is no minimal access ideal for  $2^{\aleph_0}$ .

On the other hand, the class V is V-accessible *independently* of the underlying model, in view of the axiom of regularity and the induced V-preordering  $x \preccurlyeq_{\tau} y \Leftrightarrow rank(x) \leq rank(y)$  of the universe.<sup>2</sup> Similarly, for every *regular* cardinal  $\kappa, \kappa$  is provably  $[\kappa]^{<\kappa}$ -accessible. More important are the facts shown below, that the ideals V and  $[\kappa]^{<\kappa}$  are both *provably* (a) *minimal* access ideals, and (b) the *unique* minimal access ideals for V and any regular  $\kappa$ , respectively.

Needles to say that for V, as a well as for any cardinal  $\kappa$ , there is an abundance of access ideals that arise in the form of  $Seg(\preccurlyeq)$  for the various preorderings  $\preccurlyeq$  of V or  $\kappa$ , so the proof of existence of *minimal* such ideals and, further, of *unique minimal* ones is by no means trivial. For example partition  $\kappa$  into finitely or infinitely many pieces and then rearrange the pieces into an arbitrary total ordering in which each piece keeps its standard ordering. Then a total ordering of  $\kappa$  arises and hence an access ideal for it.

According to Definition 2.1, if a class X is  $\mathcal{I}$ -accessible for an ideal  $\mathcal{I}$  within a model  $\mathfrak{M}$ , then there is a preordering  $\preccurlyeq$  of X in  $\mathfrak{M}$  such that every

 $<sup>^{2}</sup>$ Actually this fact is the reason that in the definition of accessibility 2.1 given above we use preorderings rather than total orderings or well-orderings.

 $\preccurlyeq$ -bounded set of X belongs to  $\mathcal{I}$ , i.e.,

$$\mathfrak{M} \models (\forall Y \subseteq X)(Y \text{ is } \preccurlyeq \text{-bounded} \to Y \in \mathcal{I}).$$

$$(1)$$

This definition obviously allows the ideal  $\mathcal{I}$  to contain also unbounded subsets of X. If, however,  $\mathcal{I}$  is minimal, then the converse of (1) holds too. Namely the following holds.

**Lemma 2.2** Let  $\mathfrak{M} \models \text{GBC}$  and let  $\mathcal{I}$  be a minimal access ideal for a definable class X in  $\mathfrak{M}$ . Then for every  $\mathcal{I}$ -preordering  $\preccurlyeq \in \mathfrak{M}$  of X:

$$\mathfrak{M} \models (\forall Y \subseteq X)(Y \text{ is } \preccurlyeq \text{-bounded } \leftrightarrow Y \in \mathcal{I}).$$

*Proof.* One direction follows from Definition 2.1. For the converse, assume that there is an  $\mathcal{I}$ -preordering  $\preccurlyeq \in \mathfrak{M}$  of X such that  $\mathcal{I}$  contains a set  $Y_0 \subseteq X$  unbounded in  $\langle X, \preccurlyeq \rangle$ . Then  $\mathcal{I}$  is not a minimal access ideal for X. Because X is  $Seg(\preccurlyeq)$ -accessible, and  $Seg(\preccurlyeq) \subsetneq \mathcal{I}$ , since  $Y_0 \in \mathcal{I} \backslash Seg(\preccurlyeq)$ .  $\dashv$ 

#### Minimality and uniqueness of V as an access ideal for V

**Theorem 2.3** (GBC) The class V is a minimal access ideal for V.

Proof. By the regularity axiom of GBC, the relation  $x \preccurlyeq y$  iff  $rank(x) \le rank(y)$  is a V-preordering, thus V is V-accessible in GBC. To prove the minimality of V, let  $\mathcal{J}$  be any ideal on V such that  $\mathcal{J} \subsetneq V$ . It suffices to see that V is not  $\mathcal{J}$ -accessible. Assume the contrary, and let  $\preccurlyeq_1$  be a preordering of V such that for every  $x \in V, \preccurlyeq_{1,x} \in \mathcal{J}$ . Pick a  $z \in V \setminus \mathcal{J}$ . If z were cofinal in  $\langle V, \preccurlyeq_1 \rangle$ , then  $V = \bigcup_{x \in z} \preccurlyeq_{1,x}$ , and since for every  $x, \preccurlyeq_{1,x} \in \mathcal{J} \subseteq V$ , by Replacement and Union  $\bigcup_{x \in z} \preccurlyeq_{1,x} = V$  would be a set, a contradiction. If, on the other hand, z were not cofinal in  $\langle V, \preccurlyeq_1 \rangle$ , we should have  $z \subseteq \preccurlyeq_{1,x}$  for some x, so  $z \in \mathcal{J}$  since  $\preccurlyeq_{1,x} \in \mathcal{J}$ , a contradiction again.

In fact V is not just a minimal access ideal for V, but it is the *unique* minimal access ideal. We shall need, first, the following simple consequence of the axiom of choice AC. This is rather folklore but we include its proof here for the reader's convenience.

**Lemma 2.4** (GBC) Let X be a proper class. Then for every cardinal  $\kappa$  there is set  $b \subset X$  such that  $|b| = \kappa$ .

Proof. Given the proper class X, let  $R = \{rank(x) : x \in X\}$ . R is a proper class of ordinals and for every such class there is in GB a strictly increasing enumeration  $R = \{\alpha_{\xi} : \xi \in On\}$ , i.e., a bijection  $F : On \to R$ such that  $\xi < \zeta \Leftrightarrow F(\xi) < F(z)$ . (This can be shown as follows: let S be the class of set functions f such that  $dom(f) = \xi \in On, rng(f) \subset R$  and for every  $\delta \in dom(f), f(\delta) = \min(R \setminus \{f(\gamma) : \gamma < \delta\})$ . S is defined by a normal formula so it is a class of GB. Then  $F : On \to R$  is defined as follows:  $F(\xi) = \alpha \Leftrightarrow (\exists f \in S)(\xi \in dom(f) \land f(\xi) = \alpha)$ .) If  $d_{\xi} = X \cap (V_{\alpha_{\xi}+1} \setminus V_{\alpha_{\xi}})$ , then  $\{d_{\xi} : \xi \in On\}$ , is a proper class of nonempty disjoint sets. Let s = $\{d_{\xi} : \xi < \kappa\}$ . s is a set of nonempty sets, so by AC it has a choice function f. If b = f[s], then b is as required.

#### **Theorem 2.5** (GBC) V is the unique minimal access ideal for V.

*Proof.* Towards reaching a contradiction assume that  $\mathcal{I} \neq V$  is another minimal access ideal for V. Then we cannot have  $V \subseteq \mathcal{I}$ , for since  $\mathcal{I} \neq V$ , that would mean that  $V \subsetneq \mathcal{I}$ , which contradicts the minimality of  $\mathcal{I}$ . Thus  $V \not\subseteq \mathcal{I}$ . Pick a set  $a \in V \setminus \mathcal{I}$  and fix an  $\mathcal{I}$ -preordering  $\preccurlyeq$  of V. (Remember that all preorderings used throughout have no last element.) It suffices to construct a preordering  $\preccurlyeq_1$  of V such that  $Seg(\preccurlyeq_1) \subsetneq \mathcal{I}$ , contradicting thus the minimality of  $\mathcal{I}$ .

Notice first that the set  $a \in V \setminus \mathcal{I}$  must be cofinal in  $\langle V, \preccurlyeq \rangle$ , otherwise  $a \subseteq \preccurlyeq_x$  for some x, and hence  $a \in \mathcal{I}$  since  $\preccurlyeq_x \in \mathcal{I}$ . Also for some  $x \in a$ ,  $\preccurlyeq_x$  must be a proper class, otherwise  $V = \bigcup_{x \in a} \preccurlyeq_x$  would be a set. Fix an element  $x_0 \in a$  such that  $\preccurlyeq_{x_0}$  is a proper class and let  $X = \preccurlyeq_{x_0}$ . Let also  $Y = V \setminus X$  be the final segment of  $(V, \preccurlyeq)$  determined by  $x_0$ . Then we can write

$$V = X \uplus Y \tag{2}$$

for the ordered disjoint union of the initial and final segments to which  $x_0$ splits  $\langle V, \preccurlyeq \rangle$ . We shall construct a preordering  $\preccurlyeq_1$  such that if  $J = Seg(\preccurlyeq_1)$ , then  $J \subsetneq \mathcal{I}$ . Let us set  $a^* = a \cap Y$ . Since a is cofinal in  $\preccurlyeq$  and infinite (because  $\preccurlyeq$  does not have a last element), so is  $a^*$ . Let  $|a^*| = \kappa$  and let  $a^* = \{y_{\xi} : \xi < \kappa\}$ . Using Lemma 2.4, pick a subset b of the proper class Xsuch that  $|b| = |a^*| = \kappa$  and let  $b = \{z_{\xi} : \xi < \kappa\}$ . The idea is to cut b off X and make it cofinal in the new preordering  $\preccurlyeq_1$ , retaining  $\preccurlyeq$  to the rest parts of V. This can be obtained if we define  $\preccurlyeq_1$  so that  $\preccurlyeq_1$  agrees with  $\preccurlyeq$  on  $V \setminus b$ , while b and  $a^*$  both run cofinally to V (and to each other) with respect to  $\preccurlyeq_1$ . So let

$$R = [\preccurlyeq \cap (V \setminus b)^2] \cup \{ \langle y_{\xi}, z_{\xi} \rangle, \langle z_{\xi}, y_{\xi} \rangle : \xi < \kappa \},\$$

and let  $\preccurlyeq_1$  be the transitive closure of the binary relation R, i.e.,  $\langle x, z \rangle \in \preccurlyeq_1$ whenever there is y such that  $\langle x, y \rangle$  and  $\langle y, z \rangle$  belong to  $\preccurlyeq_1$ . By this definition  $y_{\xi} \preccurlyeq_1 z_{\xi}$  and  $z_{\xi} \preccurlyeq_1 y_{\xi}$  for all  $\xi < \kappa$ , and  $\preccurlyeq_1$  is a preordering of V which, in analogy to (2), splits V into the ordered union

$$V = (X \setminus b) \uplus (Y \cup b), \tag{3}$$

with  $X \setminus b$ , Y and b being pairwise disjoint. Moreover  $a^* \subseteq Y$ , hence  $a^* \cap b = \emptyset$ and  $a^*$  is cofinal in  $\langle V, \preccurlyeq_1 \rangle$ . Let

$$\mathcal{J} = Seg(\preccurlyeq_1) = \{Z : (\exists w) (Z \subseteq \preccurlyeq_{1,w})\}.$$

We show first that  $\mathcal{J} \subseteq \mathcal{I}$ . Let  $Z \in \mathcal{J}$  and let  $Z \subseteq \preccurlyeq_{1,w}$  for some w. We shall show that  $Z \in \mathcal{I}$ . By the picture of V in (3) as an ordered union, Z is written as a disjoint union  $Z = X_1 \cup X_2 \cup x_3$ , where  $X_1 \subseteq X \setminus b, X_2 \subseteq Y$  and  $x_3 \subseteq b$ . Since  $X \setminus b$  and b already belong to  $\mathcal{I}$ , it follows that  $X_1, x_3 \in \mathcal{I}$ . It remains to show that  $X_2 \in \mathcal{I}$ . We have  $X_2 \subseteq Y$  and  $X_2 \subseteq \preccurlyeq_{1,w}$ . Clearly either  $w \in Y$  or  $w \in b$ . If  $w \in Y$ , then since  $\preccurlyeq$  and  $\preccurlyeq_1$  agree on Y, we have  $X_2 \subseteq \preccurlyeq_w$ , so  $X_2 \in \mathcal{I}$  because  $\preccurlyeq$  is an  $\mathcal{I}$ -preordering. Assume  $w \in b$ , so  $w = z_{\xi}$  for some  $\xi < \kappa$ . By definition  $z_{\xi} \preccurlyeq_1 y_{\xi}$ , where  $y_{\xi} \in a^* \subseteq Y$ , i.e.,  $y_{\xi} \in Y$ . Then  $X_2 \subseteq \preccurlyeq_{1,z_{\xi}} \subseteq \preccurlyeq_{1,y_{\xi}}$ . Since  $y_{\xi} \in Y$  as before  $X_2 \subseteq \preccurlyeq_{y_{\xi}}$ , whence again  $X_2 \in \mathcal{I}$ . Thus all three pieces  $X_1, X_2, x_3$  of Z belong to  $\mathcal{I}$ , and hence  $Z \in \mathcal{I}$ . This proves  $\mathcal{J} \subseteq \mathcal{I}$ . On the other hand  $b \notin \mathcal{J}$  because it is cofinal in  $\preccurlyeq_1$ , while  $b \in \mathcal{I}$  as a subset of X which belongs to  $\mathcal{I}$ . Therefore  $\mathcal{J} \subseteq \mathcal{I}$ .

#### Minimality and uniqueness of access ideals on cardinals

Before considering access ideals on powersets, i.e., on cardinals  $2^{\kappa}$ , we must consider access ideals on general uncountable cardinals. We shall see that there is a sharp difference between regular and singular cardinals with respect to this matter. We deal first with regular cardinals.

**Theorem 2.6** (ZFC) Let  $\kappa$  be an uncountable cardinal and < be its natural well-ordering. Then:

(a)  $Seg(<) = [\kappa]^{<\kappa}$  iff  $\kappa$  is regular.

(b) If  $\kappa$  is regular, then  $[\kappa]^{<\kappa}$  is a minimal access ideal for  $\kappa$ .

(c) Moreover if  $\kappa$  is regular,  $[\kappa]^{<\kappa}$  is the unique minimal access ideal for  $\kappa$ .

*Proof.* (a) Firstly, for every  $\kappa$ ,  $Seg(<) \subseteq [\kappa]^{<\kappa}$ . Suppose  $\kappa$  is regular and  $x \in [\kappa]^{<\kappa}$ . Then x is bounded in  $\langle \kappa, < \rangle$ , thus  $x \in Seg(<)$ , so  $[\kappa]^{<\kappa} \subseteq$ 

Seg(<) and therefore  $Seg(<) = [\kappa]^{<\kappa}$ . Conversely, if  $\kappa$  is not regular there is  $x \subset \kappa$  cofinal in  $\langle \kappa, < \rangle$  with  $|x| < \kappa$ . Thus  $x \in [\kappa]^{<\kappa} \setminus Seg(<)$  and hence  $Seg(<) \subsetneq [\kappa]^{<\kappa}$ .

(b) Let  $\kappa$  be regular and assume  $[\kappa]^{<\kappa}$  is not minimal for  $\kappa$ , i.e., there is an ideal  $\mathcal{I} \subsetneq [\kappa]^{<\kappa}$  and an  $\mathcal{I}$ -preordering  $\preccurlyeq$  for  $\kappa$ . Pick  $x \in [\kappa]^{<\kappa} \setminus \mathcal{I}$ . Then x is cofinal in  $\langle \kappa, \preccurlyeq \rangle$ , because otherwise  $x \subseteq \preccurlyeq_{\xi} \in \mathcal{I}$ , for some  $\xi \in \kappa$ , and hence  $x \in \mathcal{I}$ , contradiction. This implies that  $\kappa = \bigcup_{\xi \in x} \preccurlyeq_{\xi}$ . But for each  $\xi$ ,  $| \preccurlyeq_{\xi} | < \kappa$ , since  $\preccurlyeq_{\xi} \in \mathcal{I} \subseteq [\kappa]^{<\kappa}$ , and also  $|x| < \kappa$ . This clearly contradicts the regularity for  $\kappa$ .

(c) Let  $\kappa$  be regular. The proof that  $[\kappa]^{<\kappa}$  is the unique minimal ideal for  $\kappa$  is quite similar to the proof of Theorem 2.5 for the corresponding fact about V, so we simply sketch it. Let  $\mathcal{I} \neq [\kappa]^{<\kappa}$  be another minimal access ideal for  $\kappa$  and fix an  $\mathcal{I}$ -preordering  $\preccurlyeq$  of  $\kappa$ . If  $[\kappa]^{<\kappa} \subseteq \mathcal{I}$ , then  $[\kappa]^{<\kappa} \subseteq \mathcal{I}$ , so  $\mathcal{I}$  would not be minimal. Therefore  $[\kappa]^{<\kappa} \not\subseteq \mathcal{I}$ . Pick  $a \in [\kappa]^{<\kappa} \backslash \mathcal{I}$ . Then, as in the proof of 2.5, a is cofinal in  $\langle \kappa, \preccurlyeq \rangle$ , and for some  $x \in a$ ,  $|\preccurlyeq_x| = \kappa$ . Because  $\kappa = \bigcup_{x \in a} \preccurlyeq_x$ , and  $|a| < \kappa$ , so if  $|\preccurlyeq_x| < \kappa$  for all  $x \in \kappa, \kappa$  would be singular. Fix such a  $x_0 \in \kappa$  and let  $X = \preccurlyeq_{x_0}$ , so  $|X| = \kappa, Y = \kappa \backslash X$ , and  $a^* = a \cap Y$ . Let  $|a^*| = \lambda < \kappa$ , and pick  $b \subset X$ , with  $|b| = |a^*| = \lambda$ . Let also  $a^* = \{y_{\xi} : \xi < \lambda\}$  and  $|b| = \{z_{\xi} : \xi < \lambda\}$ . We define a preordering  $\preccurlyeq_1$  of  $\kappa$  in complete analogy to that of 2.5, making b cofinal in  $\langle \kappa, \preccurlyeq_1 \rangle$ , and so that  $\kappa$ is written

$$\kappa = (X \setminus b) \uplus (Y \cup b),$$

with  $X \setminus b$ , Y and b being pairwise disjoint. If  $\mathcal{J} = Seg(\preccurlyeq_1)$ , it is shown as before that  $\mathcal{J} \subsetneq \mathcal{I}$ , which contradicts the minimality of  $\mathcal{I}$ .

Now when coming to powersets, Theorem 2.6 (c) gives its place to the following variant that claims consistency rather than provability.

**Theorem 2.7** (ZFC) For any  $\kappa \geq \aleph_0$  and any regular  $\lambda > \kappa$ , it is (relatively) consistent to believe that  $[2^{\kappa}]^{<\lambda}$  is the unique minimal access ideal on  $2^{\kappa}$ .

*Proof.* It is well-known that the following is consistent with ZFC:  $\lambda > \kappa$  is regular and  $2^{\kappa} = \lambda$ . This is obtained by a standard Cohen forcing extension  $\mathfrak{M}$  in which  $\lambda = 2^{\kappa}$  is still regular, so it follows from 2.6 (c) that in  $\mathfrak{M}$ ,  $[2^{\kappa}]^{<\lambda} = [\lambda]^{<\lambda}$  is the unique minimal access ideal on  $2^{\kappa}$ .

We come next to the  $\mathcal{I}$ -accessibility of singular cardinals. First recall that for every preordered set  $\langle A, \preccurlyeq \rangle$ , the *cofinality* of  $\langle A, \preccurlyeq \rangle$ , denoted  $\operatorname{cof}(A, \preccurlyeq)$ , is the least cardinal  $\lambda$  such that: There is a set  $X \subseteq A$  such that  $|X| = \lambda$ , and  $\langle X, \preccurlyeq \rangle$  is well-ordered and cofinal to  $\langle A, \preccurlyeq \rangle$ . Clearly, for every infinite set preordering  $\langle A, \preccurlyeq \rangle$  having no last element,  $\operatorname{cof}(A, \preccurlyeq)$  is an infinite regular cardinal. Therefore, as an immediate consequence we have the following.

**Fact 2.8** If  $\kappa$  is a singular cardinal, then for every preordering  $\preccurlyeq$  of  $\kappa$ , there is a regular  $\lambda < \kappa$  such that  $\operatorname{cof}(\kappa, \preccurlyeq) = \lambda$ .

Note that for any  $\kappa$ , either regular or singular,  $\operatorname{cof}(\kappa, \preccurlyeq)$  need not be equal to  $\operatorname{cf}(\kappa)$ .  $\operatorname{cof}(\kappa, \preccurlyeq)$  depends heavily on  $\preccurlyeq$  and may be less or greater than  $\operatorname{cf}(\kappa)$ . For example for every  $\kappa$  there is  $\preccurlyeq$  such that  $\operatorname{cof}(\kappa, \preccurlyeq) = \aleph_0$ . On the other hand  $\operatorname{cf}(\aleph_{\omega_1}) = \aleph_1$ , but we can easily find orderings of  $\aleph_{\omega_1}$ with cofinalities  $\aleph_n > \aleph_1$ . More generally, given any regular  $\lambda$  and  $\kappa > \lambda$ , we can find a preordering  $\preccurlyeq_{\lambda}$  of  $\kappa$  such that  $\operatorname{cof}(\kappa, \preccurlyeq_{\lambda}) = \lambda$ . (This can be done by splitting  $\kappa$  into  $\lambda$  (disjoint) pieces  $x_i$ ,  $i < \lambda$ , e.g. by the help of a surjection  $f : \kappa \to \lambda$ , and then define  $\preccurlyeq_{\lambda}$  by totally ordering the pieces  $x_i$ so that  $x_i \prec_{\lambda} x_j$  iff i < j.)

Next we show that, even when  $\kappa$  is singular, for every  $\lambda < \kappa$ ,  $\kappa$  is not  $[\kappa]^{<\lambda}$ -accessible, i.e., there is no  $[\kappa]^{<\lambda}$ -preordering of  $\kappa$ . This is due to the following fact.

**Lemma 2.9** Let  $\kappa$  be an uncountable cardinal and let  $\aleph_0 \leq \lambda < \kappa$ . For every preordering  $\preccurlyeq$  of  $\kappa$ , there is an  $x \in \kappa$  such that  $|\preccurlyeq_x| \geq \lambda$ .

*Proof.* Assume for contradiction that for a given  $\lambda < \kappa$ , there is a preordering  $\preccurlyeq$  of  $\kappa$  such that  $|\preccurlyeq_x| < \lambda$  for every  $x \in \kappa$ . For each  $x \in \kappa$ , let  $[x] = \{y \in \kappa : x \preccurlyeq y \& y \preccurlyeq x\}$  and let  $\kappa/\sim = \{[x] : x \in \kappa\}$ . Then  $\preccurlyeq$  is (essentially) a linear ordering of  $\kappa/\sim$ . Since  $|\preccurlyeq_x| < \lambda$  for all x, we have that  $|[x]| < \lambda$  for all  $x \in \kappa$ , and therefore  $|\kappa/\sim| = \kappa$ . So it suffices to prove the claim for the linearly ordered set  $\langle \kappa/\sim, \preccurlyeq \rangle$ . For simplicity we pick an element  $s_x$  from each class [x], and we set  $S = \{s_x : x \in \kappa\}$ . Then  $|S| = \kappa$ , and we prove the claim for the isomorphic set  $\langle S, \preccurlyeq \rangle$ . Now either S contains a cofinal subset Y of cardinality  $\lambda$ , or not. In the first case  $S = \bigcup_{s \in Y} \preccurlyeq_s$ , where  $|Y| = \lambda$  and  $|\preccurlyeq_s| < \lambda$ , for every  $s \in Y$ . But then  $|S| = \lambda$ , a contradiction. In the other case, S does not contain any cofinal set of cardinality  $\lambda$ . So if we pick a  $Y \subseteq S$  with  $|Y| = \lambda$ , then there is an  $s \in S$  such that  $Y \subseteq \preccurlyeq_s$ . But then  $|\preccurlyeq_s| \geq \lambda$ , a contradiction again.

**Corollary 2.10** Let  $\aleph_0 \leq \lambda < \kappa$ . Then for every  $\mathcal{I} \subseteq [\kappa]^{<\lambda}$ ,  $\kappa$  is not  $\mathcal{I}$ -accessible, i.e., there is no  $\mathcal{I}$ -preordering of  $\kappa$ .

*Proof.* If  $\preccurlyeq$  were a  $[\kappa]^{<\lambda}$ -preordering of  $\kappa$ , then  $|\preccurlyeq_x| < \lambda$  for every  $x \in \kappa$ , which contradicts 2.9.

Next theorem is in sharp contrast to Theorem 2.6 (c).

**Theorem 2.11** (ZFC) If  $\kappa$  is a singular cardinal, then there is no minimal access ideal on  $\kappa$ . In particular, Seg(<) is not minimal for  $\kappa$ .

*Proof.* Let  $\kappa$  be singular. We have to show that for every preordering  $\preccurlyeq$  of  $\kappa$ ,  $Seg(\preccurlyeq)$  is not minimal. That is, there is a preordering  $\preccurlyeq_1$  of  $\kappa$  such that  $Seg(\preccurlyeq_1) \subsetneq Seg(\preccurlyeq)$ . It is remarkable that the construction of  $\preccurlyeq_1$  follows exactly the same steps as the corresponding constructions in the proofs of theorems 2.5 and 2.6 (c), although the latter theorems make claims to the opposite direction from that of the present one.

Fix a preordering  $\preccurlyeq$  of  $\kappa$  and let  $\lambda = \operatorname{cof}(\kappa, \preccurlyeq)$ . By Fact 2.8,  $\lambda < \kappa$ . Also by Lemma 2.9 there is  $x_0 \in \kappa$  with  $| \preccurlyeq_{x_0} | \geq \lambda$ . We construct  $\preccurlyeq_1$  imitating the proofs of 2.5 and 2.6 (c). Namely, let  $X = \preccurlyeq_{x_0}$ , let  $b \subseteq X$  with  $|b| = \lambda$ , and let  $Y = \kappa \setminus X$ . Pick also a set a with  $|a| = \lambda$ , which is well-ordered by  $\preccurlyeq$  with order-type  $\lambda$ , and  $\langle a, \preccurlyeq \rangle$  is cofinal to  $\langle \kappa, \preccurlyeq \rangle$ . So we can write  $a = \{\alpha_{\xi} : \xi < \lambda\}$ , where  $\alpha_{\xi} \prec \alpha_{\zeta} \Leftrightarrow \xi < \zeta$ . Without loss of generality we may assume that  $a \subseteq Y$ . Let  $b = \{\beta_{\xi} : \xi < \lambda\}$ . As in the previous proofs we cut off b from X and make it cofinal to a. Namely, we let

$$R = [\preccurlyeq \cap (\kappa \backslash b)^2] \cup \{ \langle \alpha_{\xi}, \beta_{\xi} \rangle, \langle \beta_{\xi}, \alpha_{\xi} \rangle : \xi < \lambda \}.$$

Also we set  $\preccurlyeq_1 =$  the transitive closure of the binary relation R. Then  $\preccurlyeq_1$  is as required. For exactly as in the proof of 2.5 we see that  $Seg(\preccurlyeq_1) \subseteq Seg(\preccurlyeq)$ , while on the other hand  $b \in Seg(\preccurlyeq) \setminus Seg(\preccurlyeq_1)$ . Thus  $Seg(\preccurlyeq_1) \subsetneq Seg(\preccurlyeq)$  and this completes the proof.

As a complement to Theorem 2.7 we have the following.

**Theorem 2.12** (ZFC) For any cardinal  $\kappa \geq \aleph_0$ , it is consistent to believe that there is no minimal access ideal for  $2^{\kappa}$ .

*Proof.* Given  $\kappa \geq \aleph_0$ , it is well-known that the following is consistent with ZFC:  $\kappa < \operatorname{cf}(\lambda) < \lambda$  and  $2^{\kappa} = \lambda$ . Fixing any singular cardinal  $\lambda$  with  $\kappa < \operatorname{cf}(\lambda) < \lambda$ , if  $\mathfrak{M}$  is a model of ZFC obtained by a  $\kappa$ -cc forcing which makes  $2^{\kappa} = \lambda$ , then  $\lambda$  remains singular in  $\mathfrak{M}$ . So it follows from Theorem 2.11 that in  $\mathfrak{M}$  no minimal access ideal exists for  $2^{\kappa}$ .

## 3 Algebraic similarity of $\mathcal{P}(\kappa)$ and $\mathcal{C}(V)$ equipped with ideals

Given infinite cardinals  $\lambda \leq \kappa$ , let  $\langle \mathcal{P}(\kappa), [\kappa]^{<\lambda} \rangle$  and  $\langle \mathcal{C}(V), V \rangle$  be the Boolean algebras of all subcollections of  $\kappa$  and V, equipped with the ideals  $[\kappa]^{<\lambda}$  and V, respectively, representing some "small sets/classes" of the algebras. The language of the structures is  $L_b(S) = L_b \cup \{S\}$ , where  $L_b$  is the language of Boolean algebras, and  $S(\cdot)$  is a unary predicate symbol for the elements of the ideals. The purpose of this section is to prove the following.

**Theorem 3.1** (GBC) For all infinite cardinals  $\lambda \leq \kappa$ ,  $\langle \mathcal{P}(\kappa), [\kappa]^{<\lambda} \rangle \equiv_{L_b(S)} \langle \mathcal{C}(V), V \rangle$ . In particular for  $\kappa = 2^{\lambda}$  we get  $\langle \mathcal{P}(2^{\lambda}), [2^{\lambda}]^{\leq \lambda} \rangle \equiv_{L_b(S)} \langle \mathcal{C}(V), V \rangle$ .

It is already known that any two powerset Boolean algebras are elementarily equivalent. This is an easy corollary of a general result concerning algebras with the same "invariants" (see [3, p. 296], §18). However the algebras here are endowed with ideals so a stronger argument is needed. The method of proof is an application of the Ehrenfeucht-Fraïssé games (see e.g. [4], pp. 52f). We recall briefly some definitions.

Let L be a first-order language and let  $\mathcal{M} = \langle M, \ldots \rangle$ ,  $\mathcal{N} = \langle N, \ldots \rangle$  be two L-structures. We define first what an infinite game is between  $\mathcal{M}$ ,  $\mathcal{N}$ . This is a two-player game denoted  $G_{\omega}(\mathcal{M}, \mathcal{N})$ . The players, called I and II, play at stages  $1, 2, \ldots$ , by choosing at each stage i, alternately elements  $a_i \in M$  and  $b_i \in N$ . A play consists of  $\omega$  stages. Player I may choose an element either from  $\mathcal{M}$  or from N, and then II must respond by choosing an element from the other structure. Thus if I chooses  $a_i \in \mathcal{M}$  (resp.  $b_i \in N$ ), then II must choose  $b_i \in N$  (resp.  $a_i \in \mathcal{M}$ ). Player II wins the play if  $f = \{\langle a_i, b_i \rangle : i = 1, 2, \ldots\}$  is the graph of a partial embedding from  $\mathcal{M}$  into  $\mathcal{N}$ . Otherwise I wins. A strategy for player II in  $G_{\omega}(\mathcal{M}, \mathcal{N})$  is a function  $\tau$ such that for each i, if  $u_1, \ldots, u_i$  are the first i moves of player I, then the i-th move of player II is  $\tau(u_1, \ldots, u_i)$ . The strategy  $\tau$  is a winning strategy for II, if for any sequence  $u_1, u_2, \ldots$  of moves of player I, II wins the play if he/she follows the strategy  $\tau$ .

Now an Ehrenfeucht-Fraïssé (E-F) game between  $\mathcal{M}$ ,  $\mathcal{N}$  is a game of the above general form, except that it consists only of *finitely* many stages instead of infinite ones. It is denoted  $G_n(\mathcal{M}, \mathcal{N})$  if it stops after the *n*-th stage. The winning strategy for player II or I is defined as before. E-F games can be used to characterize elementary equivalence of two structures. Specifically the following holds. **Theorem 3.2** ([4, Thm 2.4.6]) Let L be a finite language without function symbols and let  $\mathcal{M}$ ,  $\mathcal{N}$  be L-structures. Then  $\mathcal{M} \equiv \mathcal{N}$  if and only if player II has a winning strategy in the game  $G_n(\mathcal{M}, \mathcal{N})$ , for all n > 0.

We shall use 3.2 to prove Theorem 3.1. In the proof we shall need the following simple fact which is provable in GB.

**Lemma 3.3** (GB) Every proper class can be partitioned into two proper disjoint subclasses.

*Proof.* Let X be a proper class and, as in Lemma 2.4, consider the class  $R = \{rank(x) : x \in X\}$  which is proper and can be written  $R = \{\alpha_{\xi} : \xi \in On\}$  in its natural well-ordering. Let  $R_0 = \{\alpha_{\xi} : \xi \text{ is a limit ordinal}\}$ , and  $R_1 = \{\alpha_{\xi} : \xi \text{ is a successor ordinal}\}$ . If for i = 0, 1 we set  $X_i = \{x \in X : rank(x) \in R_i\}, X_0, X_1$  is the required partition.

Proof of Theorem 3.1. Fix infinite cardinals  $\lambda \leq \kappa$  and let  $\mathcal{M} =$  $\langle \mathcal{P}(\kappa), [\kappa]^{<\lambda} \rangle$  and  $\mathcal{N} = \langle \mathcal{C}(V), V \rangle$ . (Remember that we can always think we are working in a model  $\mathfrak{K} = \langle K, \mathcal{K} \rangle$  of GBC, with  $V^{\mathfrak{K}} = K$  and  $\mathcal{C}(V)^{\mathfrak{K}} = \mathcal{K}$ . so in this case  $\mathcal{N} = \langle \mathcal{K}, K \rangle$ .) In order to comply with the condition of Theorem 3.2 that the language of the structures must be relational, we shall take the language  $L_b$  of Boolean algebras to be  $L_b = \{\sqsubseteq\}$  (instead of the more often used algebraic language  $\{+, \cdot, -, 0, 1\}$ ). According to Theorem 3.2, we have to show that for each n, Player II has a winning strategy in  $G_n(\mathcal{M}, \mathcal{N})$ . In fact we can show that player II has a winning strategy not only in the finite games  $G_n(\mathcal{M}, \mathcal{N})$ , but also in the infinite game  $G_{\omega}(\mathcal{M}, \mathcal{N})$ . To see this it suffices to define, for every n, what the move of player II will be after the completion of the first n stages. Now since our ground structures  $\mathcal{M}, \mathcal{N}$ are Boolean algebras, every finite subset A of each of the algebras generates a finite Boolean subalgebra *fully determined* by the set of its atoms which are the intersections of the elements of A. And whenever a player chooses a new element from the ground algebra, this element expands the previous subalgebra to a larger finite one by refining the set of its atoms. Thus in our case each stage of the game is not obtained from the previous one by the attachment of a single new pair of elements from each algebra, but of several new pairs that correspond to the generated new atoms. So we can consider the rounds of the game as partial embeddings  $f_i : C_i \to D_i, i \ge 1$ , where  $C_i, D_i$  are finite Boolean subalgebras of  $\mathcal{P}(\kappa)$  and  $\mathcal{C}(V)$  respectively, and for every  $i, f_i \subsetneq f_{i+1}$ . Moreover,  $f_i$  respect the predicate  $S(\cdot)$  by preserving the ideals, i.e., satisfying the condition

(\*) For every 
$$a \in C_i$$
,  $a \in [\kappa]^{<\lambda} \Leftrightarrow f_i(a) \in V$ .

Assume that the first n rounds of the game have been completed, i.e., embeddings  $f_i : C_i \to D_i$  which satisfy (\*) have been defined for all  $i \leq n$ , and suppose player I starts the (n + 1)-th round by choosing either some  $a \in \mathcal{P}(\kappa)$  or some  $b \in \mathcal{C}(V)$ . We have to define the response of player II, i.e.  $a \ b \in \mathcal{C}(V)$ , if player I chooses  $a \in \mathcal{P}(\kappa)$ , or an  $a \in \mathcal{P}(\kappa)$ , if player I chooses  $b \in \mathcal{C}(V)$ , and the extension  $f_{n+1} : C_{n+1} \to D_{n+1}$  of  $f_n$  so that  $\langle a, b \rangle \in f_{n+1}$ and  $f_{n+1}$  satisfies (\*). Without loss of generality we assume that  $a \in \mathcal{P}(\kappa)$ . We shall describe the definition of b and  $f_{n+1}$ , and this will be the description of the winning strategy of player II. A quite similar strategy would be described if player I had started with a  $b \in \mathcal{C}(V)$  and player II should choose  $a \in \mathcal{P}(\kappa)$ .

The strategy of player II: Let  $c_1, \ldots, c_k$  be the atoms of the algebra  $C_n$ which form a partition of  $\kappa$ . Then  $f_n(c_1), \ldots, f_n(c_k)$  are the atoms of  $D_n$ which form a partition of V. For each  $i = 1, \ldots, k$ , let  $x_i = c_i \cap a$  and  $x'_i = c_i \cap (\kappa \setminus a)$  and let

$$X = \{x_i, x'_i \neq \emptyset : i = 1, \dots, k\}.$$

X is a new partition of  $\kappa$  which refines the partition of  $c_i$ 's and determines a finite subalgebra  $C_{n+1}$  of  $\mathcal{P}(\kappa)$  such that  $a \in C_{n+1} \supseteq C_n$ . In order to define  $f_{n+1}$  and  $f_{n+1}(a) = b$ , clearly it suffices to define  $f_{n+1}(x_i) = y_i$  and  $f_{n+1}(x'_i) = y'_i$  for all  $x_i, x'_i \in X$ , so that  $y_i, y'_i$  form a suitable partition of  $f_n(c_i)$ . In particular we must have

$$|x_i| < \lambda \ (|x_i'| < \lambda) \Leftrightarrow y_i \in V \ (y_i' \in V).$$

$$\tag{4}$$

If for some  $i, c_i \subseteq a$ , or  $c_i \subseteq \kappa \setminus a$ , then  $x_i = c_i$  or  $x'_i = c_i$ , respectively, for which  $f_n$  is already defined. So if  $x_i = c_i$  or  $x'_i = c_i$ , we just let  $f_{n+1}(x_i) = f_n(c_i)$  and  $f_{n+1}(x'_i) = f_n(c_i)$ . In all other cases  $x_i, x'_i$  form a proper partition of  $c_i$  and we consider the following cases with respect to the cardinalities of  $x_i$  and  $x'_i$ .

1)  $|x_i|, |x'_i| \ge \lambda$ . Then  $|c_i| \ge \lambda$  so, by (\*) for  $f_n, f_n(c_i)$  is a proper class. In this case it suffices to take, using Lemma 3.3,  $y_i, y'_i$  to be complementary proper subclasses of  $f_n(c_i)$ .

2)  $|x_i| < \lambda$  and  $|x'_i| \ge \lambda$  (or vice versa). Then  $|c_i| \ge \lambda$ , hence  $f_n(c_i)$  is a proper class. Choose  $\emptyset \ne y_i \subseteq f_n(c_i)$  to be a set and let  $y'_i = f_n(c_i) \setminus y_i$ . In particular, if  $x_i$  is finite, we pick  $y_i$  so that  $|y_i| = |x_i|$ .

3)  $|x_i|, |x'_i| < \lambda$ . Then  $|c_i| < \lambda$ . Therefore  $f_n(c_i)$  is a set. If  $x_i, x'_i$  are both infinite, we pick any infinite  $y_i, y'_i$  that form a partition of  $f_n(c_i)$ . If

either  $x_i$  or/and  $x'_i$  is finite we take the partition  $y_i, y'_i$  so that  $|y_i| = |x_i|$ or/and  $|y'_i| = |x'_i|$ , respectively. (This is always possible, even when  $c_i$  is finite. Because the same strategy has been followed by player II at the previous stages, so if  $c_i$  is finite the strategy has yielded  $f_n(c_i)$  such that  $|f_n(c_i)| = |c_i|$ .) Let

$$Y = \{y_i, y'_i : i = 1, \dots, k\}.$$

Y is the set of atoms of a subalgebra  $D_{n+1} \supseteq D_n$  of  $\mathcal{C}(V)$  and the function  $f_{n+1}: C_{n+1} \to D_{n+1}$  for which  $f_{n+1}(x_i) = y_i$  and  $f_{n+1}(x'_i) = y'_i$  induces a partial isomorphism that extends  $f_n$ . Thus clearly  $f_{n+1}$  satisfies (4). It satisfies also condition (\*) above. Because by (4),  $f_{n+1}$  already satisfies (\*) for the atoms of  $C_{n+1}$ . If now x is any element of  $C_{n+1}$  and  $z_1, \ldots, z_m$  are the atoms of  $C_{n+1}$  below x, then  $x = z_1 \cup \cdots \cup z_m$ , hence  $f_{n+1}(x) = f_{n+1}(z_1) \cup \cdots \cup f(z_m)$ . So clearly  $|x| < \lambda$  iff  $|z_i| < \lambda$  for all  $i = 1, \ldots, m$ , iff  $f_{n+1}(z_i)$  is a set for all  $i = 1, \ldots, m$ , iff  $f_{n+1}(x)$  is a set. Thus  $f_{n+1}$  satisfies (\*).

This completes the description of the strategy of player II. It remains to prove that the strategy is winning, that is, all these  $f_i$  are partial embeddings from  $\mathcal{M}$  into  $\mathcal{N}$ . Namely it suffices to show the following:

Claim. Let F be the set of all finite functions from  $\mathcal{M}$  into  $\mathcal{N}$  which are defined through E-F games, in which player II follows the above described strategy. Then every  $f \in F$  is a partial embedding, i.e., for every  $\phi(x_1, \ldots, x_n)$  of  $L_b(S)$  and every  $a_1, \ldots, a_n \in \mathcal{P}(\kappa)$  such that  $a_1, \ldots, a_n \in$ dom(f),

$$\mathcal{M} \models \phi(a_1, \dots, a_n) \Leftrightarrow \mathcal{N} \models \phi(f(a_1), \dots, f(a_n)).$$
(5)

Proof of Claim. We prove that (5) holds for every  $f \in F$  and for every  $\phi$  by induction on the length of  $\phi$ . The atomic formulas of  $L_b(S)$  are  $x \sqsubseteq y$  and S(x) and clearly the ordering  $\sqsubseteq$  is preserved by all  $f \in F$ . Also, we saw above that every  $f \in F$  satisfies condition (\*), so for every  $a \in dom(f)$ 

$$\mathcal{M} \models S(a) \Leftrightarrow a \in [\kappa]^{<\lambda} \Leftrightarrow f(a) \in V \Leftrightarrow \mathcal{N} \models S(f(a)).$$

The induction steps for the connectives are straightforward for all  $f \in F$ , and the only nontrivial step of the induction is the quantifier step. So given a  $\phi(x, x_1, \ldots, x_n)$  we assume that all  $f \in F$  satisfy (5) for  $\phi$ , and prove this for  $\exists x \phi(x, x_1, \ldots, x_n)$ , i.e., for all  $a_1, \ldots, a_n$  in dom(f)

$$\mathcal{M} \models \exists x \phi(x, a_1, \dots, a_n) \Leftrightarrow \mathcal{N} \models \exists x \phi(x, f(a_1), \dots, f(a_n)).$$
(6)

Suppose  $\mathcal{M} \models \exists x \phi(x, a_1, \dots, a_n)$  for given  $a_1, \dots, a_n$  in dom(f). Then  $\mathcal{M} \models \phi(a, a_1, \dots, a_n)$  for some  $a \in \mathcal{P}(\kappa)$ . Now, by definition, f is the *n*-th

stage of some E-F game between  $\mathcal{M}$  and  $\mathcal{N}$ , so a can be the move of player I at the (n + 1)-th stage of the game. Then player II responds following the above strategy by picking a  $b \in \mathcal{C}(V)$  and extending f to  $g \supseteq f \cup \{\langle a, b \rangle\}$  so that  $g \in F$ . By our induction assumption g satisfies (5) for  $\phi$ , hence  $\mathcal{N} \models \phi(g(a), g(a_1), \ldots, g(a_n))$ , or  $\mathcal{N} \models \phi(b, f(a_1), \ldots, f(a_n))$ , therefore  $\mathcal{N} \models \exists x \phi(x, f(a_1), \ldots, f(a_n))$ . This proves " $\Rightarrow$ " of (6). For the converse assume  $\mathcal{N} \models \exists x \phi(x, f(a_1), \ldots, f(a_n))$  and  $\mathcal{N} \models \phi(b, f(a_1), \ldots, f(a_n))$  for a  $b \in \mathcal{C}(V)$ . As before b can also be the move of player I at the (n+1)-th stage, and then player II responds by choosing an  $a \in \mathcal{P}(\kappa)$  and  $g \supseteq f \cup \{\langle a, b \rangle\}$  so that g belongs to F. Then  $\mathcal{N} \models \phi(g(a), g(a_1), \ldots, g(a_n))$  and by the induction assumption  $\mathcal{M} \models \phi(a, a_1, \ldots, a_n)$ , whence  $\mathcal{M} \models \exists x \phi(a_1, \ldots, a_n)$ . This proves the Claim and completes the proof of the Theorem.

The similarity of structures  $\langle \mathcal{P}(\kappa), [\kappa]^{<\lambda} \rangle$  and  $\langle \mathcal{C}(V), V \rangle$  expressed by Theorem 3.1 is of a rather algebraic character, since it is about Boolean algebras equipped with ideals. However it can be transcribed into one with more set-theoretic flavor, i.e., one using the language of set theory. Specifically, the elements of the ground set  $\kappa$  and class V can be treated as *atoms* (or *urelements*), while  $\mathcal{P}(\kappa)$  and  $\mathcal{C}(V)$  as collections of sets/classes of atoms. So let  $L_1(S)$  be the language consisting of  $\in$ ,  $S(\cdot)$ , variables  $p, q, \ldots$  for urelements, and variables  $x, y, \ldots$  for sets. The atomic formulas of  $L_1(S)$  are:  $p = q, p \in x$ , and S(x). The intended  $L_1(S)$ -structures are  $\langle \kappa, \mathcal{P}(\kappa), [\kappa]^{<\lambda} \rangle$ , for sets, and  $\langle V, \mathcal{C}(V), V \rangle$  for the universe. However a technical problem here is the fact that the two kinds of entities, atoms and sets, should be disjoint, while  $\kappa \cap \mathcal{P}(\kappa) \neq \emptyset$  (and similarly for V). We can bypass this obstacle by working with suitable copies of  $\kappa$  and V. Namely we set

$$\kappa^* = \{ \langle x, 0 \rangle : x \in \kappa \}$$
 and  $V^* = \{ \langle x, 0 \rangle : x \in V \}.$ 

Then  $\kappa^* \cap \mathcal{P}(\kappa^*) = \emptyset$  and  $V^* \cap \mathcal{C}(V^*) = \emptyset$ . Under this coding the  $L_b(S)$  structures  $\langle \mathcal{P}(\kappa), [\kappa]^{<\lambda} \rangle$  and  $\langle \mathcal{C}(V), V \rangle$  become isomorphic to  $\langle \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle$  and  $\langle \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle$ , respectively, where  $\mathcal{P}(V^*)$  is the class of subsets of  $V^*$ . Then in view of Theorem 3.1 we have: For all infinite cardinals  $\lambda \leq \kappa$ 

$$\langle \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle \equiv_{L_b(S)} \langle \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle.$$
(7)

Based on (7) we can easily deduce the following.

**Theorem 3.4** For all infinite  $\lambda \leq \kappa$ ,

$$\langle \kappa^*, \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle \equiv_{L_1(S)} \langle V^*, \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle.$$

*Proof.* Define an interpretation  $\phi \mapsto \phi^+$  from the atomic formulas of  $L_1(S)$  to formulas of  $L_b(S)$  as follows: Fix enumerations  $p_n, x_n, n \in \mathbb{N}$ , of the variables of  $L_1(S)$  for atoms and sets, respectively, and an enumeration  $y_n, n \in \mathbb{N}$ , of the variables of  $L_b(S)$ . The atomic formulas of  $L_1(S)$  are  $p_i = p_j, p_i \in x_j$  and  $S(x_j)$ . Define

- (a)  $(p_i = p_j)^+ := (y_{2i} = y_{2j}),$ (b)  $(p_i \in x_j)^+ := Atom(y_{2i}) \land y_{2i} \sqsubseteq y_{2j+1},$
- (c)  $\tilde{S}(x_j)^+ = S(y_{2j+1}),$

where Atom(y) is the formula of  $L_b(S)$  saying that "y is an atom". This interpretation is also extended to all formulas of  $L_1(S)$  as expected. (As follows from clauses (a)-(c) above, the variable  $p_i$  translates to  $y_{2i}$  while  $x_i$ translates to  $y_{2j+1}$ . This translation determines also of course the interpretation of quantifiers.) Then it is easy to see that for all  $\phi \in L_1(S)$ ,

$$\langle \kappa^*, \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle \models \phi \iff \langle \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle \models \phi^+$$

and

$$\langle V^*, \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle \models \phi \Leftrightarrow \langle \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle \models \phi^+$$

By (7),

$$\langle \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle \models \phi^+ \Leftrightarrow \langle \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle \models \phi^+,$$

therefore

$$\langle \kappa^*, \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle \models \phi \iff \langle V^*, \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle \models \phi.$$

In fact Theorems 3.1 and 3.4 are essentially equivalent, since a proof of 3.1 from 3.4 can also be obtained by defining a converse interpretation  $\phi \mapsto \phi'$  from atomic formulas of  $L_b(S)$  to formulas of  $L_1(S)$  as follows:

$$(x \sqsubseteq y)' := \forall p(p \in x \to p \in y),$$
  

$$(x = y)' := (x = y), \text{ and}$$
  

$$S(x)' := S(x).$$

This interpretation, extended to all formulas, yields: For every  $\phi$  of  $L_b(S)$ ,

$$\langle \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle \models \phi \iff \langle \kappa^*, \mathcal{P}(\kappa^*), [\kappa^*]^{<\lambda} \rangle \models \phi$$

and

$$\langle \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle \models \phi \Leftrightarrow \langle V^*, \mathcal{C}(V^*), \mathcal{P}(V^*) \rangle \models \phi'.$$

Thus 3.4 implies 3.1.

## 4 Order-theoretic similarity of uncountable cardinals with the class of ordinals

By the axiom of choice every uncountable set A can be identified with an uncountable cardinal  $\kappa$ , and by the axiom of global choice vN the universe Vcan be identified with the class of ordinals On. However these identifications add structure to A and V, namely the well-orderings of  $\kappa$  and On, which we denote by the same symbol <. Admitting the naturalness of choice principles, it is then reasonable to compare the first-order structures  $\langle \kappa, < \rangle$ ,  $\langle \lambda, < \rangle$  for uncountable  $\kappa, \lambda$ , as well as  $\langle \kappa, < \rangle$  and  $\langle On, < \rangle$ . In this section we shall show how the main result of [7] (Theorem 2.1), which says that the monadic  $\forall_1^1$  positive theory of linear ordering is the same for all uncountable cardinals, can be extended also over  $\langle On, < \rangle$ .

The starting point of [7] was the following result, which is a special case of [2, Cor. 44]:

**Theorem 4.1** ([7, Thm 1.1]) For all uncountable cardinals  $\kappa$ ,  $\lambda \langle \kappa, \langle \rangle \equiv_{L_{ord}} \langle \lambda, \langle \rangle$ , where  $L_{ord} = \{ \prec \}$  is the first-order language of ordering.

Remember that the ordering < of ordinals is a well-ordering, that is a secondorder property, while the equivalence  $\equiv_{L_{ord}}$  refers only to the first-order properties of <. Theorem 4.1 is in fact a consequence of Corollary 44 of [2] and, essentially, the same argument that we used in order to justify it in [7] can be applied also to justify the following extension.

**Theorem 4.2** For every uncountable cardinal  $\kappa$ ,  $\langle \kappa, \langle \rangle \equiv_{L_{ord}} \langle On, \langle \rangle$ .

*Proof.* (Sketch) By Corollary 44 of [2], for any ordinals  $\alpha$ ,  $\beta$ ,  $\langle \alpha, \langle \rangle \equiv \langle \beta, \langle \rangle$  if and only if  $\alpha$ ,  $\beta$  are congruent modulo  $\omega^{\omega}$ , i.e., there are  $\xi$ ,  $\eta$ ,  $\delta$  such that  $\delta < \omega^{\omega}$ ,  $\alpha = \omega^{\omega} \cdot \xi + \delta$ ,  $\beta = \omega^{\omega} \cdot \eta + \delta$  and either  $\xi = \eta = 0$  or  $\xi \neq 0$  and  $\eta \neq 0.^3$  Since for all cardinals  $\kappa, \lambda > \omega$ ,  $\kappa = \omega^{\omega} \cdot \kappa$  and  $\lambda = \omega^{\omega} \cdot \lambda$ , it follows that  $\kappa, \lambda$  are congruent modulo  $\omega^{\omega}$ . This already proves Theorem 4.1.

Concerning the justification of the present claim, without any risk of fallacy we can treat On in this particular situation as the "greatest (limit) ordinal", and extend the usual (right) multiplication of ordinals over  $On \cup \{On\}$  in the obvious way. Specifically, just as for any  $\alpha$  and any limit  $\beta$  we

<sup>&</sup>lt;sup>3</sup>Important notice: when I was preparing [7], Professor Doner had kindly informed me in a private communication, that the definition of *congruence modulo*  $\omega^{\omega}$  as given in [2], p. 51, is mistaken. The correct definition is that given above and can be found in [1, p. 6].

have  $\alpha \cdot \beta = \sup\{\alpha \cdot \gamma : \gamma < \beta\}$ , simply define  $\alpha \cdot On = \{\alpha \cdot \gamma : \gamma < On\}$ . Then trivially  $\omega^{\omega} \cdot On = On$ , so for every uncountable cardinal  $\kappa$ ,  $\kappa$  and On are congruent modulo  $\omega^{\omega}$ . Applying again Corollary 44 of [2] to  $\langle \kappa, < \rangle$  and  $\langle On, < \rangle$ , we are done. (As for the details why the relation of congruence modulo  $\omega^{\omega}$  entails the elementary equivalence of well-ordered sets or classes, one should consult [2], especially the theory  $\mathcal{W}$  (Def. 2) of "weak well-orderings", the notion of "canonical order type" (Def. 39), and Theorem 41 which relates these notions with elementarily equivalent ordinals.)  $\dashv$ 

We proceed next to extend the main theorem of [7], which establishes a similarity significantly stronger than that of Theorem 4.1, so as to hold between every uncountable cardinal and On. We need first some terminology and notation.

We first extend the language of ordering  $L_{ord} = \{\prec\}$  used previously for the structures  $\langle \kappa, < \rangle$  (with lower-case individual variables  $x, y, z, \ldots$ ) to the monadic (second-order) language  $L_{mon}$  which is  $L_{ord}$  augmented with  $\in$  and upper-case set variables  $X, Y, Z, \ldots$ . A formula of  $L_{mod}$  without quantified set variables is called normal. A formula of the form  $\forall X\phi$  (resp.  $\exists X\phi$ ), where  $\phi$  is normal is called a  $\forall_1^1$  (resp.  $\exists_1^1$ ) formula. The standard interpretations of  $\mathcal{L}_{mon}$  are the structures  $\langle A, <, \mathcal{P}(A) \rangle$  where  $\langle A, < \rangle$  is an ordered set that interprets the formulas of  $L_{ord}$ , and  $\mathcal{P}(A)$  is the range of set variables. However there are more general interpretations. They are of the form  $\langle A, <, A \rangle$ , where  $\mathcal{A} \subset \mathcal{P}(A)$ .  $\mathcal{A}$  may be some class of "small" sets, i.e., an ideal of  $\mathcal{P}(A)$ , or, in the opposite direction, a class of "large" sets. Most often these large sets are the sets cofinal in  $\langle A, < \rangle$ , so let us set

$$\operatorname{Cof}(A) = \{ X \subseteq A : X \text{ is cofinal in } \langle A, \langle \rangle \}.$$

A can be also a proper class and  $\mathcal{A}$  be a class of subclasses of A. When we deal with an ordered collection which is a proper class, like On, then, in order to interpret the set variables X, Y, we must work again in a class theory. Then as usual we write  $\mathcal{C}(On)$  for the collection of all subclasses of On. In particular below we shall deal with the analogue of  $\operatorname{Cof}(\kappa)$  above, which is

 $Cof(On) = \{ X \subseteq On : X \text{ is cofinal in } \langle On, < \rangle \}$ 

and consists of the proper subclasses of On.

**Definition 4.3** Let  $\phi(X)$  be a normal formula of  $\mathcal{L}_{mon}$  with at most one (free) set variable X.  $\phi(X)$  is said to be *positive in* X (or just *positive*) if it belongs to the smallest class of formulas which (a) contains all formulas

of  $L_{ord}$ , (b) contains the atomic formulas  $x \in X$  and (c) is closed under the logical operations  $\land$ ,  $\lor$ , and the first-order quantifiers  $\exists$  and  $\forall$ .

The main result of [7] is that the structures  $\langle \kappa, <, \operatorname{Cof}(\kappa) \rangle$ , for all uncountable cardinals  $\kappa$ , are elementarily equivalent with respect to  $\forall_1^1$  positive formulas of  $L_{mod}$ .

**Theorem 4.4** ([7, Thm. 2.1]) In ZFC the following holds: For all uncountable cardinals  $\kappa$ ,  $\lambda$ ,

$$\langle \kappa, <, \operatorname{Cof}(\kappa) \rangle \equiv_{pos}^{\forall_1^1} \langle \lambda, <, \operatorname{Cof}(\lambda) \rangle,$$

*i.e.*, for every positive  $\phi(X)$ ,

 $\langle \kappa, \langle \operatorname{Cof}(\kappa) \rangle \models (\forall X) \phi(X) \Leftrightarrow \langle \lambda, \langle \operatorname{Cof}(\lambda) \rangle \models (\forall X) \phi(X).$ 

Here we strengthen 4.4 to the following.

**Theorem 4.5** (GBC) For every uncountable cardinal  $\kappa$ ,

$$\langle \kappa, <, \operatorname{Cof}(\kappa) \rangle \equiv_{pos}^{\forall_1^1} \langle On, <, \operatorname{Cof}(On) \rangle.$$

*Proof.* (Sketch) The proof goes exactly along the lines of the proof of 4.4 and is based on the syntactic form of positive (normal) formulas of  $L_{mod}$  which has been discovered by Y. Moschovakis. Namely, every formula  $\phi(X)$  positive in X is equivalent over any structure  $\langle A, <, \mathcal{P}(A) \rangle$ , for all  $X \neq A$ , to a formula having prenex form  $(\overline{Q}\overline{w})(\forall u)(\theta(\overline{w}, u) \rightarrow u \in X)$ , where  $\overline{w} = \langle w_1, \ldots, w_n \rangle$  is a string of variables,  $\overline{Q} = \langle Q_1, \ldots, Q_n \rangle$  a string of quantifiers and  $\theta(\overline{w}, u)$  is quantifier-free. (Note that, since  $\phi(X)$  does not contain set quantifiers, the above prenex form for  $\phi(X)$  is valid also over every structure  $\langle A, \langle A \rangle$ , with  $\mathcal{A} \subseteq \mathcal{P}(A)$ .) In view of the preceding syntactic characterization of  $\forall_1^1$  positive formulas, the proof splits into several cases that correspond to the various forms of the string  $\overline{Q}\overline{w}$ , as well as to the forms of  $\theta$  as a Boolean combination of atoms and negated atoms. Then, through an exhaustive inspection of all these cases, it is shown that we can essentially reduce the truth of  $(\forall X)\phi(X)$  to the truth of a *first-order* formula. But this means, in view of Theorem 4.2, that  $(\forall X)\phi(X)$  is absolute between  $\langle \kappa, \langle \operatorname{Cof}(\kappa) \rangle$  and  $\langle On, \langle \operatorname{Cof}(On) \rangle$ .  $\dashv$ 

If we work in GB + vN, where a bijection between On and V exists, the well-ordering of On is transferred to a well-ordering of V so the following holds.

**Corollary 4.6** In GB + vN it is provable that there is a well-ordering  $\prec$  of V such that for every uncountable cardinal  $\kappa$ ,

$$\langle \kappa, \langle, \operatorname{Cof}(\kappa) \rangle \equiv_{pos}^{\forall_1^1} \langle V, \prec, \operatorname{Cof}(V) \rangle$$

Concerning the optimality of Theorem 4.5 with respect to the indistinguishability of structures  $\langle \kappa, <, \operatorname{Cof}(\kappa) \rangle$  and  $\langle On, <, \operatorname{Cof}(On) \rangle$ , and the analogous question about structures  $\langle \kappa, <, \operatorname{Cof}(\kappa) \rangle$  and  $\langle \lambda, <, \operatorname{Cof}(\lambda) \rangle$  for uncountable  $\kappa$ ,  $\lambda$ , the reader is referred to the corresponding discussion in [7]: It is not known whether  $\langle \kappa, <, \operatorname{Cof}(\kappa) \rangle$  and  $\langle \lambda, <, \operatorname{Cof}(\lambda) \rangle$  are indistinguishable with respect to all  $\forall_1^1$  and  $\exists_1^1$  formulas. What we know is that the structures  $\langle \omega_1, <, \operatorname{Cof}(\omega_1) \rangle$  and  $\langle \omega_2, <, \operatorname{Cof}(\omega_2) \rangle$  can be distinguished by a sentence  $\phi_{\omega_1}$  of the form  $(\forall x)(\exists X)\psi$ , where  $\psi$  is normal but not positive (see section 5 of [7]). This is easily generalized to all  $\omega_m, \omega_n$ . A fortiori  $\phi_{\omega_1}$ distinguishes also the structures  $\langle \omega_1, <, \operatorname{Cof}(\omega_1) \rangle$  and  $\langle On, <, \operatorname{Cof}(On) \rangle$ .

Acknowledgement My thanks go to an anonymous referee whose remarks helped me fix some unclarities and improve considerably the presentation of the article.

## References

- J.E. Doner, Definability in extended arithmetic of ordinal numbers, Dissertationes Math. (Rozprawy Mat.) 96 (1972), 50 pp.
- J.E. Doner, A. Mostowski and A. Tarski, The elementary theory of well-ordering – a metamathematical study, in: *Logic Colloquium '77*, A. Macintyre, L. Pacholski and J. Paris (eds.), North Holland, 1978, pp. 1-54.
- [3] S. Koppelberg, Metamathematics (of Boolean a lgebras), in: *Handbook of Boolean Algebras*, J.D. Monk and R. Bonnet (eds.), vol. 1, Elsevier Science Publishers B.V. 1989, pp. 287-307.
- [4] D. Marker, Model Theory: An Introduction, Springer, 2002.
- [5] R. Parikh, Existence and feasibility in arithmetic, J. Symb. Logic 36 (1971), no. 3, 494-508.
- [6] A. Tzouvaras, Positive set-operators of low complexity, Math. Log. Quar. 49 (2003), no. 3, 284-292.

- [7] A. Tzouvaras, Uncountable cardinals have the same monadic  $\forall_1^1$ -positive theory over large sets, *Fund. Math.* **181** (2004), no. 2, 125-142.
- [8] A. Tzouvaras, What is so special with the power set operation?, Arch. Math. Logic 43 (2004), no. 6, 723-737.
- [9] A. Tzouvaras, Localizing the axioms, Arch. Math. Logic 49 (2010), no. 5, 571-601.